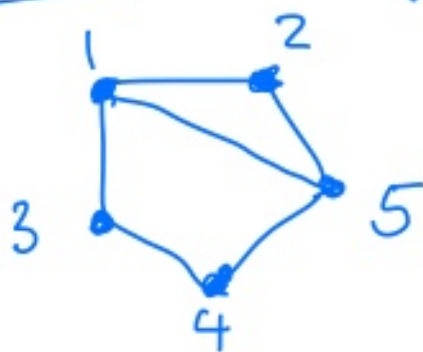


Question: In the graph below ① Find the adjacency matrix.



② How many paths go from vertex 2 to vertex 3 that have length 3?

Answers ①

$$\begin{pmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \end{pmatrix}$$

② $2 \rightarrow 5 \rightarrow 4 \rightarrow 3$ (2 ways)
 $2 \rightarrow 5 \rightarrow 1 \rightarrow 3$

Given a graph G , a path is a sequence of edges of the form

(undirected) $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\}, \dots, \{v_n, v_{n+1}\}.$

(directed) $(v_1, v_2), (v_2, v_3), \dots, (v_n, v_{n+1})$

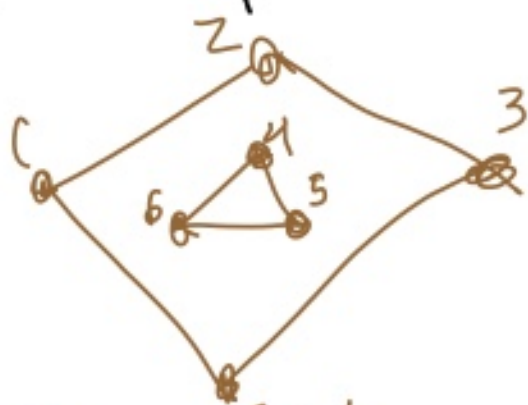
↑ These are called paths of length n , because they include n edges.

Note: We could repeat the same edge more than once, and also vertices can be used multiple times.

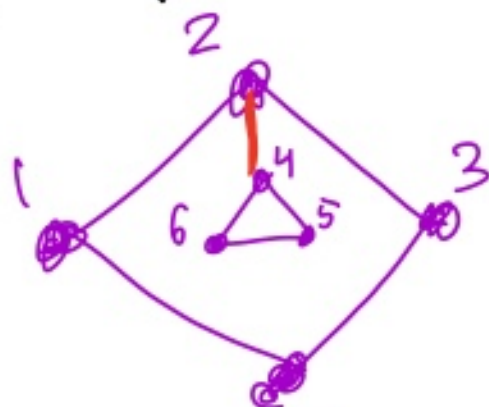
A closed path is a path that starts and ends

at the same vertex. (Another name for closed path is circuit.) A simple path is a path with no repeated edges.

A graph is called connected if we can find a path between any two vertices.



This graph is disconnected (ie. not connected), because for instance there is no path from vertex 3 to vertex 5.



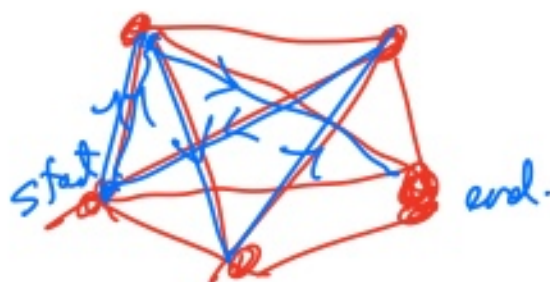
Now it's connected

Theorem: If a graph G is connected, then there exists a simple path between any two vertices.

In fact, we can connect any two vertices with a simple path where no vertices are repeated.

Proof:

Idea



Sketch of the proof: Suppose there is a path between vertices a & b in G , given by $\{a, v_2\}, \{v_2, v_3\}, \dots, \{v_n, b\}$
 $\uparrow \qquad \qquad \qquad \uparrow$
 $v_1 \qquad \qquad \qquad v_{n+1}$

If $v_j = v_k$ for some $j < k$, then
This path also works.

$\{a, v_2\}, \dots, \{v_{j-1}, v_j\}, \overbrace{\{v_k, v_{k+1}\}}^{\text{same}}, \dots, \{v_n, b\}$

Now we have 1 less repetition of vertices.

Keep shortening the path until there are no remaining repeated vertices. $\square$

Corollary: If G has N vertices — then if G is connected, then there exists a simple path between any two vertices of length $\leq N$.

Measures of how connected a graph is:

A vertex cut of a graph G is a set of vertices, that, when removed from G (including the edges attached to those vertices), the resulting graph is disconnected. An edge cut of a graph G is a set of edges such that, when they are removed, the graph is disconnected.

A couple of invariants of Graphs related to connectivity:

$\kappa(G)$ = minimum # of vertices in a vertex cut of a graph

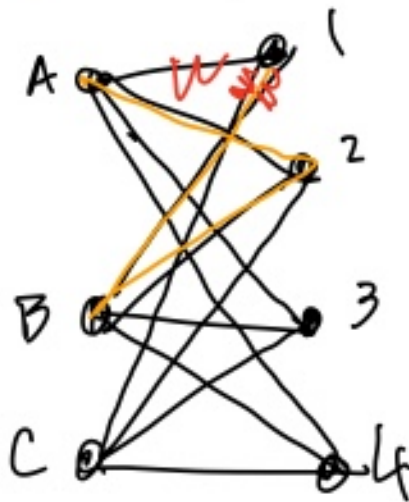
$\kappa(K_n) = n-1$

$\lambda(G)$ = minimum # of edges in an edge cut of a graph.

Thm. - For a given graph,

$$\kappa(G) \leq \lambda(G) \leq \left(\text{minimum deg of a vertex on the graph} \right).$$

Example $K_{3,4}$



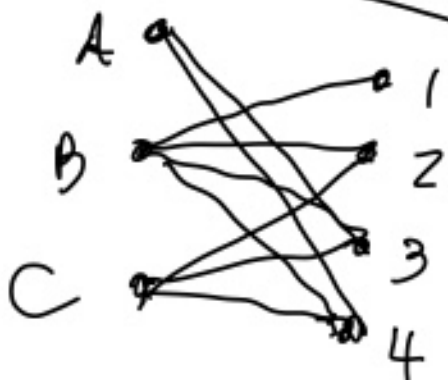
edge cut - delete $\{A,1\}, \{B,1\}, \{C,1\}$
1 disconnected from the rest

in fact $\lambda(G) = 3$

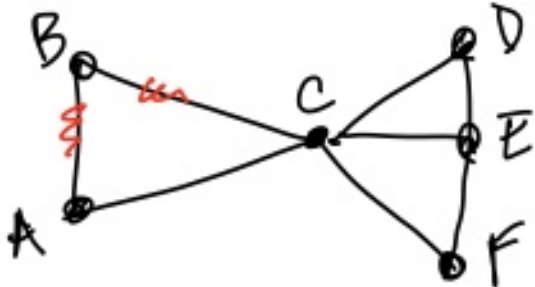
vertex cut delete A, B, C

$\kappa(G) = 3$

mindeg = 3



edge cut = $\{B,1\}$.



vertex cut : remove C.
 $K(G)=1$

$$\lambda(G)=2$$

edge cut: $\{A,B\}, \{B,C\}$.

How is this related to Adjacency matrices?

Aside: Matrix multiplication

$$\begin{pmatrix} 1 & 2 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} -1 & 2 \\ 4 & 5 \end{pmatrix} = \begin{pmatrix} 1(-1) + 2(4) & 1(2) + 2(5) \\ 0(-1) + 1(4) & 0(2) + 1(5) \end{pmatrix}$$

row column row • column

2nd row 1st column

$$= \begin{pmatrix} 7 & 12 \\ 4 & 5 \end{pmatrix}$$

$$\begin{pmatrix} \text{ } \end{pmatrix} \begin{pmatrix} \text{ } \end{pmatrix} \Rightarrow \begin{pmatrix} \text{ } \end{pmatrix}$$

$n \times \textcircled{k}$ $\textcircled{k} \times r$

row column row column

$n \times r$

Quiz

① What is your favorite ice cream flavor?

② Find the adjacency matrix of this graph

